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TAROC: TCP-Aware RObust Header Compression Scheme

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Outline



- Introduction
- Header compression/ decompression model
- Basic concepts of TAROC
- Details of TAROC
 - ✦ robustness of TAROC (window-based LSB encoding)
 - ✦ efficiency of TAROC (TCP congestion window tracking)
 - ✦ behaviors in compressor and decompressor
- Compression performance analysis

Introduction



● Related works

◆ RFC 1144 (VJHC)

- delta coding is not robust against error
- needs **TIME-OUT** to get re-synchronization

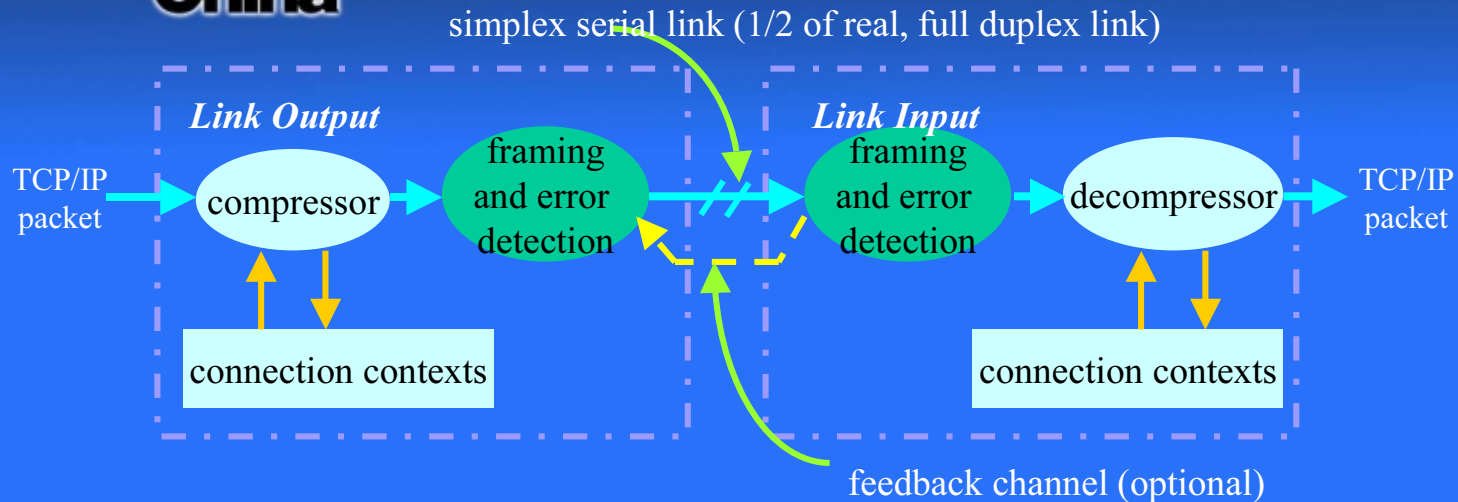
◆ RFC 2507 (IPHC)

- twice: rely on the fixed delta
 - ❖ not suited with Timestamp and SACK options
 - ❖ not suited with burst loss that may happen before decompressor
- header request: feedback channel is needed, not efficient with long RTT channel

Compression/decompression model of TAROC



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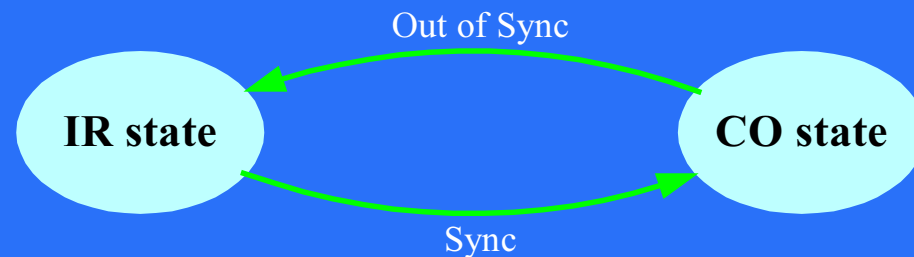
● Key ideas of TAROC

- Window-based LSB encoding
- TCP-congestion window tracking
- Slow-start, Congestion-avoidance
- Fast-retransmission, Fast-recovery

Basic concepts of TAROC (1)



- **TAROC: TCP-Aware RObust header Compression**
- Compressor starts from lower compression state and progressively transitions to higher compression state
- Compressor states:
 - ◆ IR (Initialization & Recovery):
compressor sends full headers with certain pattern to ensure the consistency of the connection context on both side
 - ◆ CO (COmpression)
higher compression state



Basic concepts of TAROC (2)



- Higher compression ratio can be achieved only if the compressor has enough confidence that the decompressor has correctly received the info
- Confidence is not obtained by acknowledgments from the decompressor
- Rather, it is obtained from the TCP congestion control window

Window-based LSB encoding (1)



- Compressor maintains values of previous packets in window VSW
- Compressor sends k least significant bits of original value
 - ◆ k is smallest value to ensure correct decompression
 - ◆ k is determined by the length of VSW
- Robustness against errors
 - ◆ compressor chooses k so that no matter what value in VSW is chosen as reference by the decompressor, decompression is correct
- The length of VSW can be shrunk
 - ◆ once compressor knows that at least one packet in VSW was received
 - ◆ no feedback channel: the length of VSW can be controlled by TCP congestion window size

Window-based LSB encoding (2)



● Reference packet selection in decompressor

- ◆ reference may be the one which is the last received non-retransmitted value, or
- ◆ uncompressed value that passes the TCP checksum successfully

● Whether a compressed or uncompressed value is sent, this value should be added into *VSW*

● Applicable fields

- ◆ TCP Sequence Number/ Acknowledge Number / Window
- ◆ IP Packet ID
- ◆ TCP Timestamp option

Tracking-based TCP congestion window estimation (1)



- Reconstruct the congestion control behavior of TCP sender at the compressor
- Adjust congestion window size in each state based on seqno

✦ Slow-Start (SS)

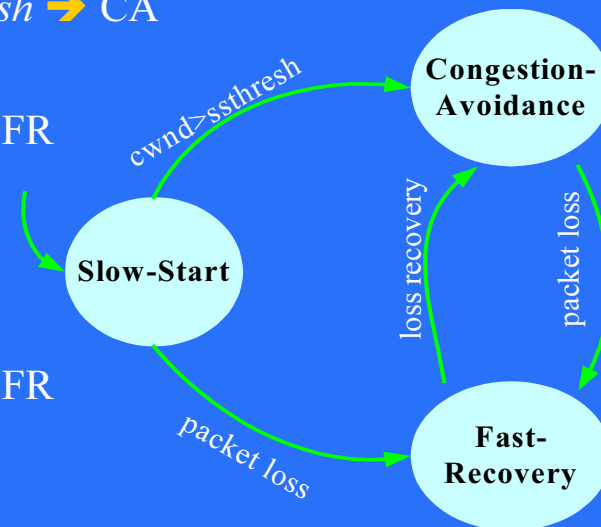
- new ordered segment: $cwnd += d$, if $cwnd > ssthresh \rightarrow CA$
- out-of-order segment: ignore it
- retransmission segment: set $cwnd$ and $ssthresh \rightarrow FR$

✦ Congestion-Avoidance (CA)

- new ordered segment: $cwnd += d * MSS / cwnd$
- out-of-order segment: ignore it
- retransmission segment: set $cwnd$ and $ssthresh \rightarrow FR$

✦ Fast-Recovery (FR)

- new segment indicating recovery of loss $\rightarrow CA$



Tracking-based TCP congestion window estimation (2)



Adjust congestion window size in each state based on ackno

Slow-Start (SS)

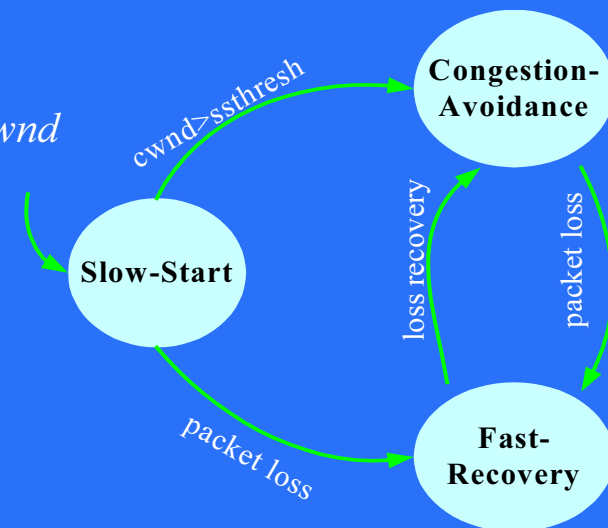
- new ordered ack: $ndupacks = 0$, $cwnd += d$, if $cwnd > ssthresh \rightarrow$ CA
- duplicate ack: $ndupacks ++$,
if $ndupacks \geq 3$, set $cwnd$ and $ssthresh \rightarrow$ FR
- otherwise, $ndupacks = 0$

Congestion-Avoidance (CA)

- new ordered ack: $ndupacks = 0$, $cwnd += d * MSS / cwnd$
- duplicate ack: $ndupacks ++$,
if $ndupacks \geq 3$, set $cwnd$ and $ssthresh \rightarrow$ FR
- otherwise, $ndupacks = 0$

Fast-Recovery (FR)

- new segment indicates recovery of loss:
 $ndupacks = 0 \rightarrow$ CA
- otherwise, $ndupacks = 0$



Compressor behavior



Initialization & Recovery (IR) State

- ✦ two types of packets in this state
 - full header packet + compressed packet
- ✦ send full header packet if
 - it satisfies slow-start pattern or it is a retransmitted packet in *VSW*
- ✦ add packet into *VSW*
- ✦ tracking congestion window
- ✦ adjust the length of *VSW* if possible

$$|VSW| \leq K \times \text{MAX}(\text{MAX}(cwnd_seq, 2 \times ssthresh_seq), \text{MAX}(cwnd_ack, 2 \times ssthresh_ack))$$

$$\text{F_PERIOD} (2^n \geq cwnd)$$

$$(1 + 2^0 + 1 + 2^1 + 1 + 2^2 + \dots + 1 + 2^{n-1} = 2^n + n - 1) > cwnd$$

Compressor behavior



● COmpression (CO) state

- ◆ only compressed packet in this state
- ◆ compress packet
- ◆ add packet into *VSW*
- ◆ tracking congestion window
- ◆ adjust the length of *VSW* if possible

$$|VSW| \leq K \times \text{MAX}(\text{MAX}(cwnd_seq, 2 \times ssthresh_seq), \\ \text{MAX}(cwnd_ack, 2 \times ssthresh_ack))$$

Decompressor behavior



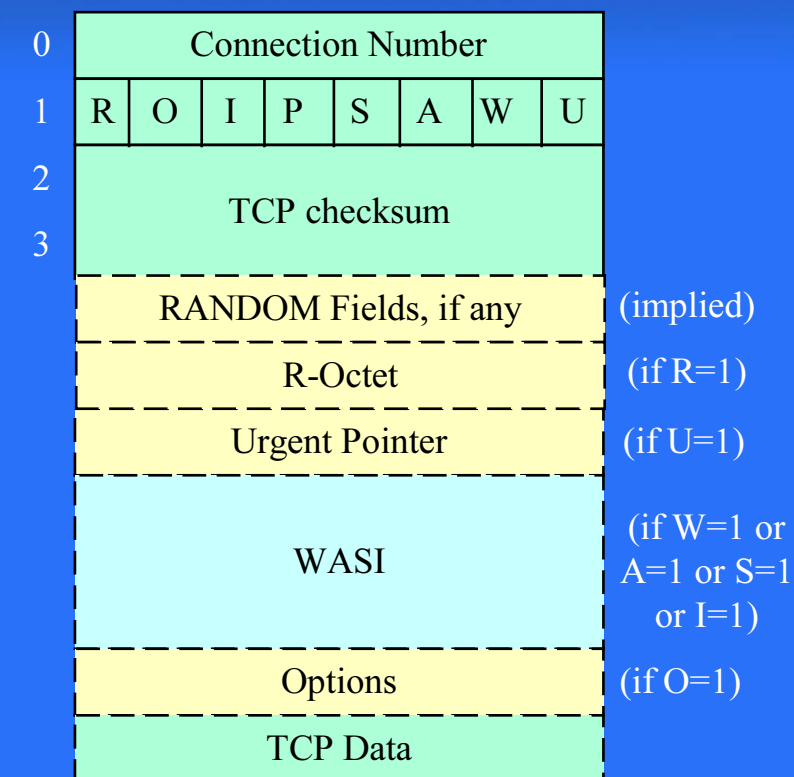
- Reference packet selection in decompressor
 - ◆ reference may be the one that is the last received non-retransmitted value, or
 - ◆ uncompressed value that passes the TCP checksum successfully
- Decompressed value determination
 - ◆ de-compressor chooses the one, which is closest to reference and whose k LSBs equal the compressed value that has been received

Compressed TCP/IP header structure



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Byte



Mandatory

Optional

Mandatory / Optional

- Almost the same as the one in IPHC except for
 - ✦ Window, Acknowledgement Number, Sequence Number, IP-ID and Options fields
- W-LSB encoding applied to
 - ✦ Window, Acknowledgement Number, Sequence Number and IP-ID
 - ✦ TS value, TS echo reply in TCP Timestamp option
- Delta coding applied to
 - ✦ SACK option

WASI field



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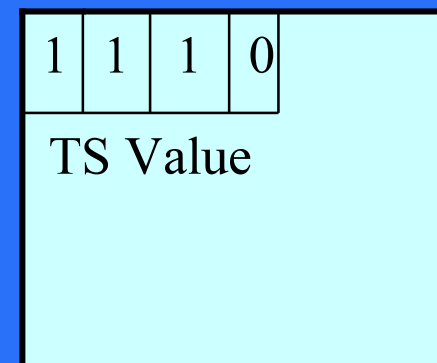
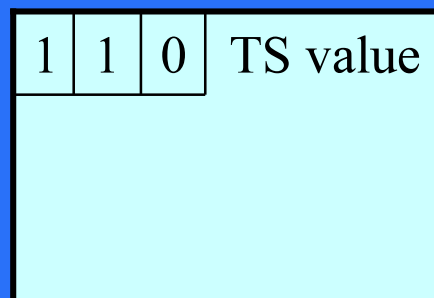
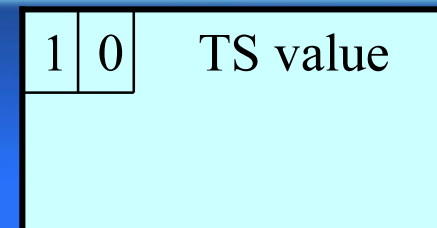
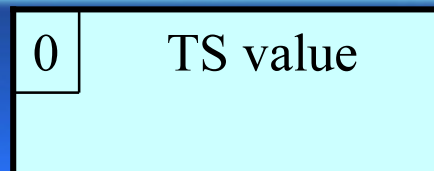
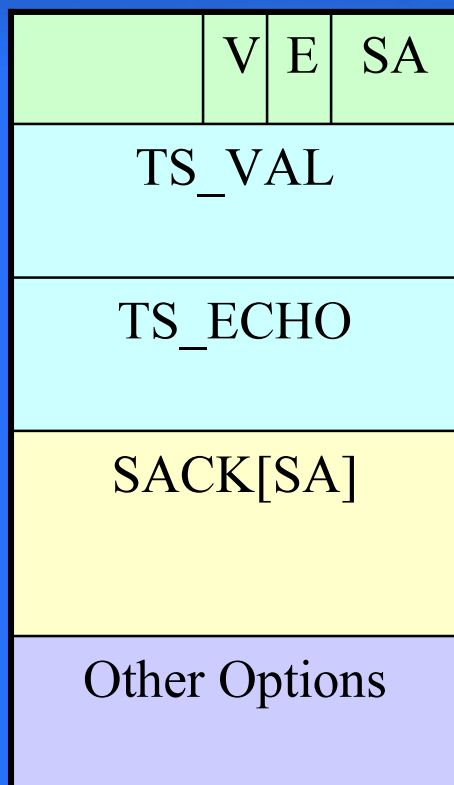
PT	C_WIN	C_ACK	C_SEQ	C_IP_ID	PADDING
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PT value	C_WIN length (bits)	C_ACK length (bits)	C_SEQ length (bits)	C_IP_ID length (bits)	length of WASI (bits)
00	10	11	11	3	2~5
01	14	15	15	7	3~7
10	16	19	19	11	4~9
11	16	23	23	15	5~11

Options field (Timestamp and SACK)

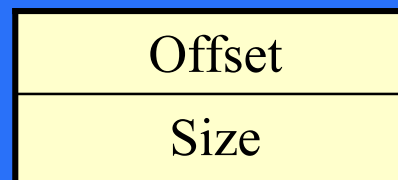


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SA: *00* no SACK option *01* 1 SACK block

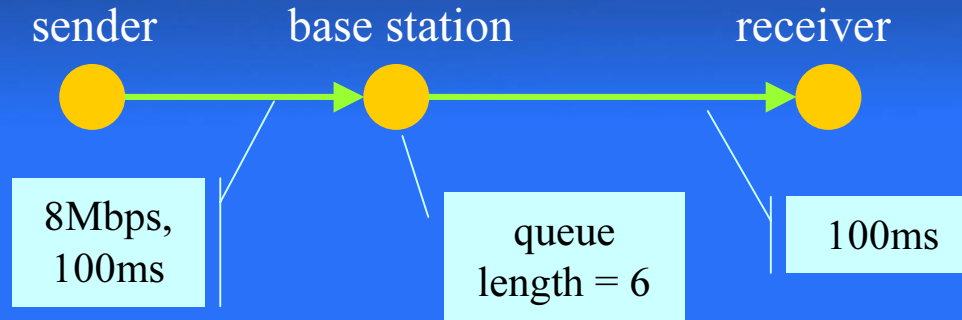
10 2 SACK blocks *11* 3 SACK blocks



Performance analysis



● Simulation topology



● Tested header compression schemes

- no header compression (None)
- RFC 1144 (VJHC)
- RFC 2507 (IPHC)
- our scheme (TAROC)
- no error happens and has minimal header size (4) (Ideal)

● Bit error rate in wireless channel (10^{-9} to 10^{-3})

● Bandwidth in wireless channel (9.6Kbps to 384Kbps)

● TCP versions: Tahoe, Reno, Sack

Performance comparison: 9.6kb/s for Sack



BER	MTU (Byte)	Performance	None	VJHC	IPHC	TAROC	Ideal
10 ⁻⁷	576	Throughput (Byte/s)	1115	1176	1175	1184	1190
		Improvement Percentage (%)	0	5.47	5.38	6.19	6.73
10 ⁻⁶	576	Throughput (Byte/s)	1112	1131	1140	1179	1185
		Improvement Percentage (%)	0	1.71	2.52	6.1	6.73
10 ⁻⁵	296	Throughput (Byte/s)	976	859	929	1100	1122
		Improvement Percentage (%)	0	-11.99	-4.82	12.7	14.96
10 ⁻⁴	168	Throughput (Byte/s)	342	218	250	459	470
		Improvement Percentage (%)	0	-36.26	-26.9	34.21	37.43
3*10 ⁻⁴	96	Throughput (Byte/s)	78	71	78	153	168
		Improvement Percentage (%)	0	-8.97	0	96.15	115.38
10 ⁻³	96	Throughput (Byte/s)	1	1	2	7	9
		Improvement Percentage (%)	0	0	0	600	800

Performance comparison: 64kb/s for Reno



BER	MTU (Byte)	Performance	None	VJHC	IPHC	TAROC	Ideal
10 ⁻⁷	576	Throughput (Byte/s)	7274	7573	7483	7642	7704
		Improvement Percentage (%)	0	4.11	2.87	5.06	5.91
10 ⁻⁶	576	Throughput (Byte/s)	6962	6577	6343	7316	7389
		Improvement Percentage (%)	0	-5.53	-8.89	5.08	6.13
10 ⁻⁵	296	Throughput (Byte/s)	3399	2461	2297	3636	3665
		Improvement Percentage (%)	0	-27.6	-32.42	6.97	7.83
10 ⁻⁴	168	Throughput (Byte/s)	455	300	322	595	607
		Improvement Percentage (%)	0	-34.07	-29.23	30.77	33.41
3*10 ⁻⁴	96	Throughput (Byte/s)	91	84	85	176	186
		Improvement Percentage (%)	0	-7.69	-6.59	93.41	104.4
10 ⁻³	96	Throughput (Byte/s)	1	1	2	9	10
		Improvement Percentage (%)	0	0	100	800	900

Performance comparison: 114kb/s for Sack



BER	MTU (Byte)	Performance	None	VJHC	IPHC	TAROC	Ideal
10 ⁻⁷	576	Throughput (Byte/s)	12004	12291	12127	12515	12557
		Improvement Percentage (%)	0	2.39	1.02	4.26	4.61
10 ⁻⁶	576	Throughput (Byte/s)	11204	9806	9160	11724	11760
		Improvement Percentage (%)	0	-12.48	-18.24	4.64	4.96
10 ⁻⁵	296	Throughput (Byte/s)	3936	2611	2599	4225	4239
		Improvement Percentage (%)	0	-33.66	-33.97	7.34	7.7
10 ⁻⁴	168	Throughput (Byte/s)	492	308	339	645	657
		Improvement Percentage (%)	0	-37.4	-31.1	31.1	33.54
3*10 ⁻⁴	96	Throughput (Byte/s)	93	84	89	183	196
		Improvement Percentage (%)	0	-9.68	-4.3	96.77	110.75
10 ⁻³	96	Throughput (Byte/s)	1	1	2	9	10
		Improvement Percentage (%)	0	0	100	800	900

Performance comparison: 384kb/s for Tahoe

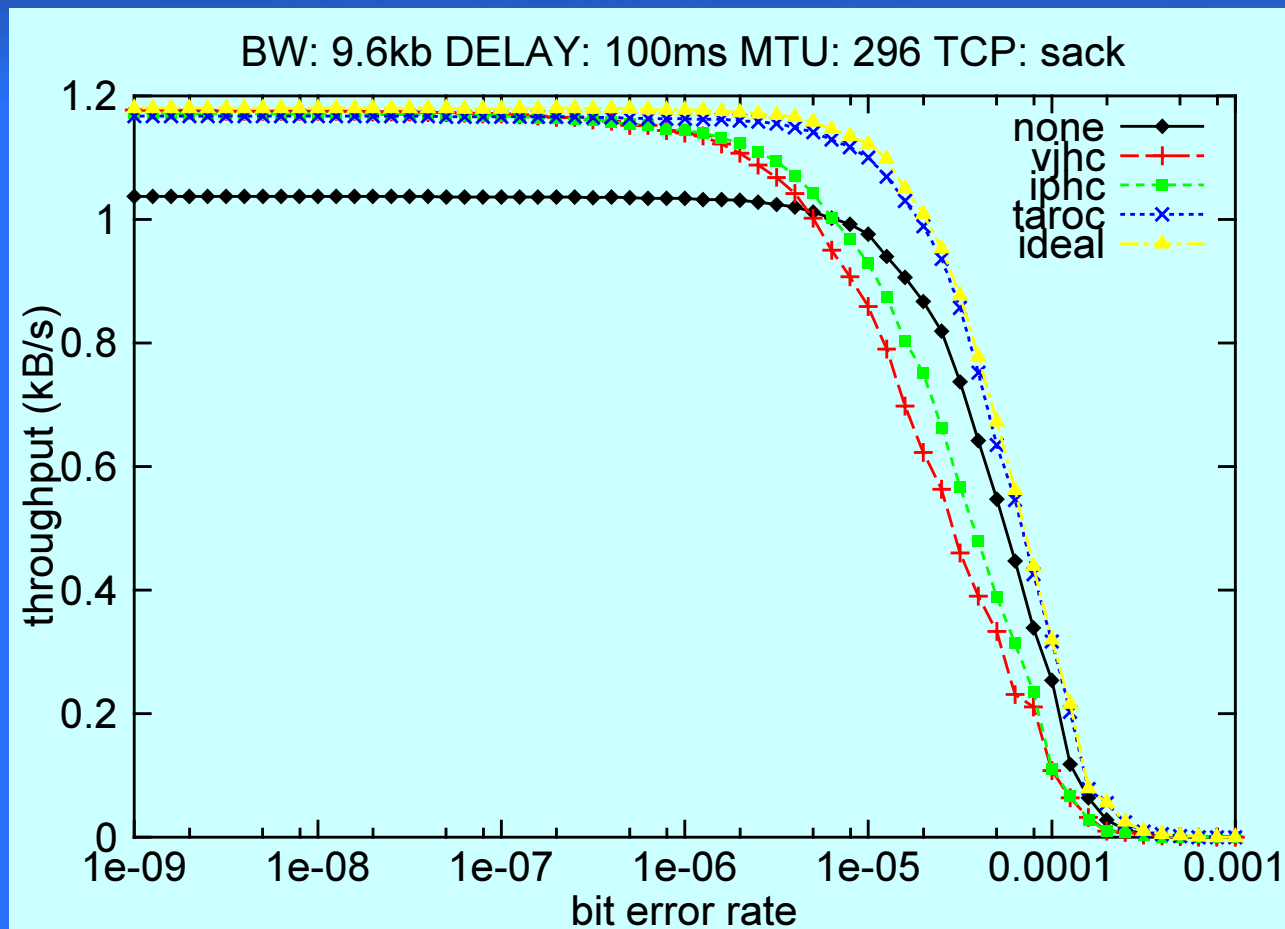


BER	MTU (Byte)	Performance	None	VJHC	IPHC	TAROC	Ideal
10 ⁻⁷	576	Throughput (Byte/s)	24650	24246	23702	24732	24753
		Improvement Percentage (%)	0	-1.64	-3.85	0.33	0.42
10 ⁻⁶	576	Throughput (Byte/s)	16710	14725	13755	17148	17183
		Improvement Percentage (%)	0	-11.88	-17.68	2.62	2.83
10 ⁻⁵	296	Throughput (Byte/s)	3548	2750	2927	3820	3835
		Improvement Percentage (%)	0	-22.49	-17.5	7.67	8.09
10 ⁻⁴	168	Throughput (Byte/s)	493	324	363	625	636
		Improvement Percentage (%)	0	-24.28	-26.37	26.77	29.01
3*10 ⁻⁴	96	Throughput (Byte/s)	92	85	89	184	195
		Improvement Percentage (%)	0	-7.61	-3.26	100	111.96
10 ⁻³	96	Throughput (Byte/s)	1	1	2	9	11
		Improvement Percentage (%)	0	0	100	800	1000

Performance comparison: 9.6kb/s for Sack



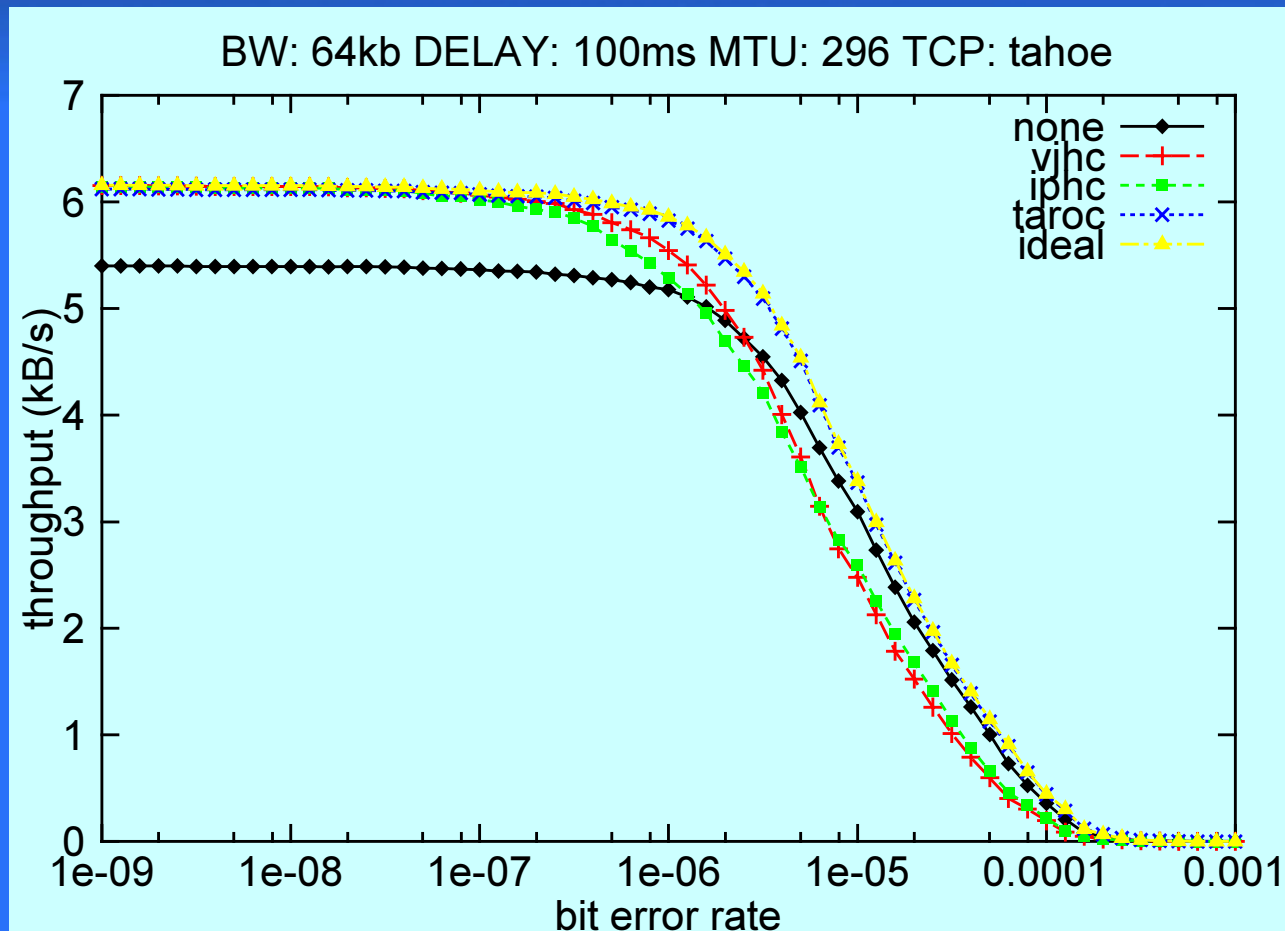
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Performance comparison: 64kb/s for Tahoe



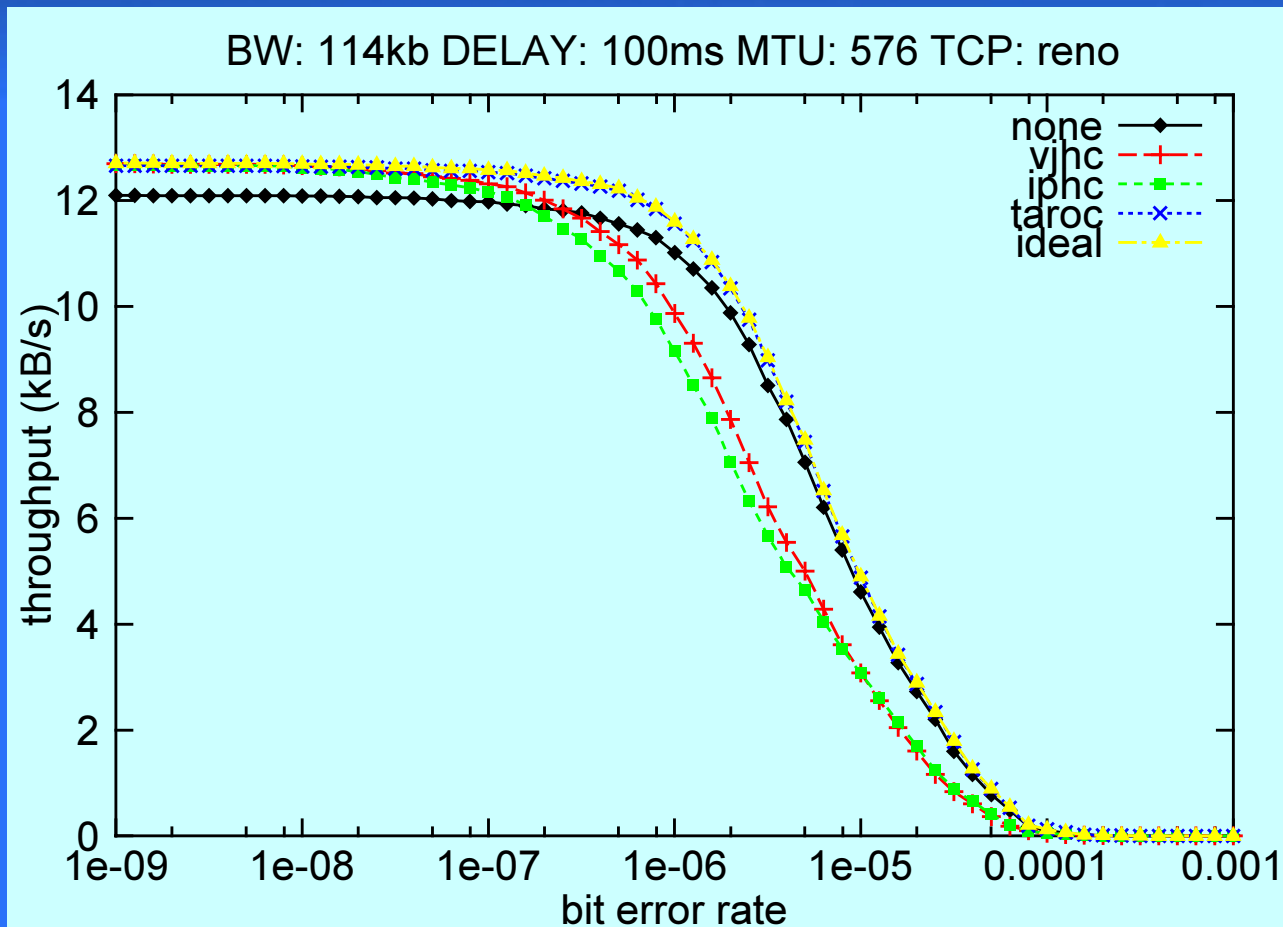
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Performance comparison: 114kb/s for Reno



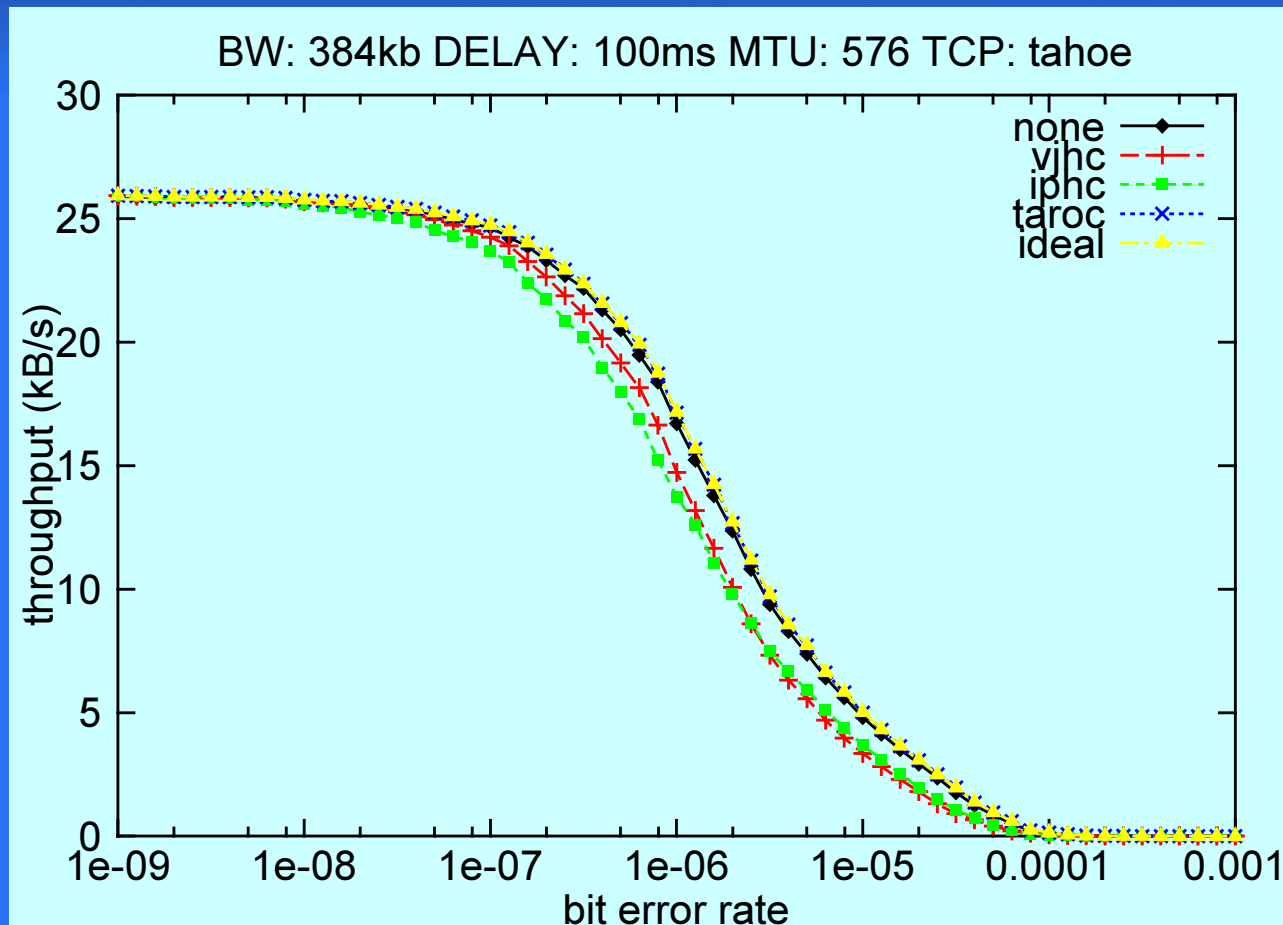
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Performance comparison: 384kb/s for Tahoe



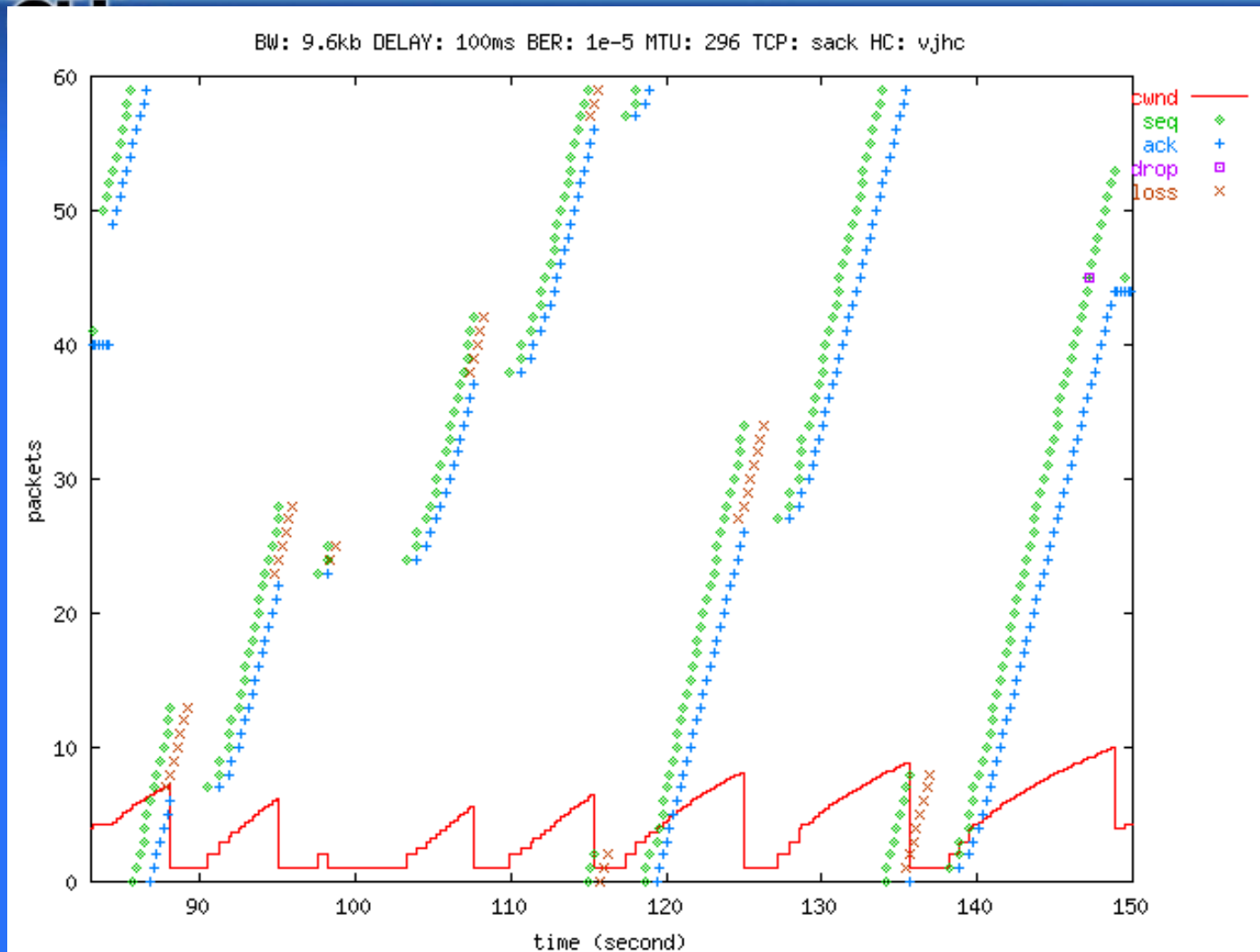
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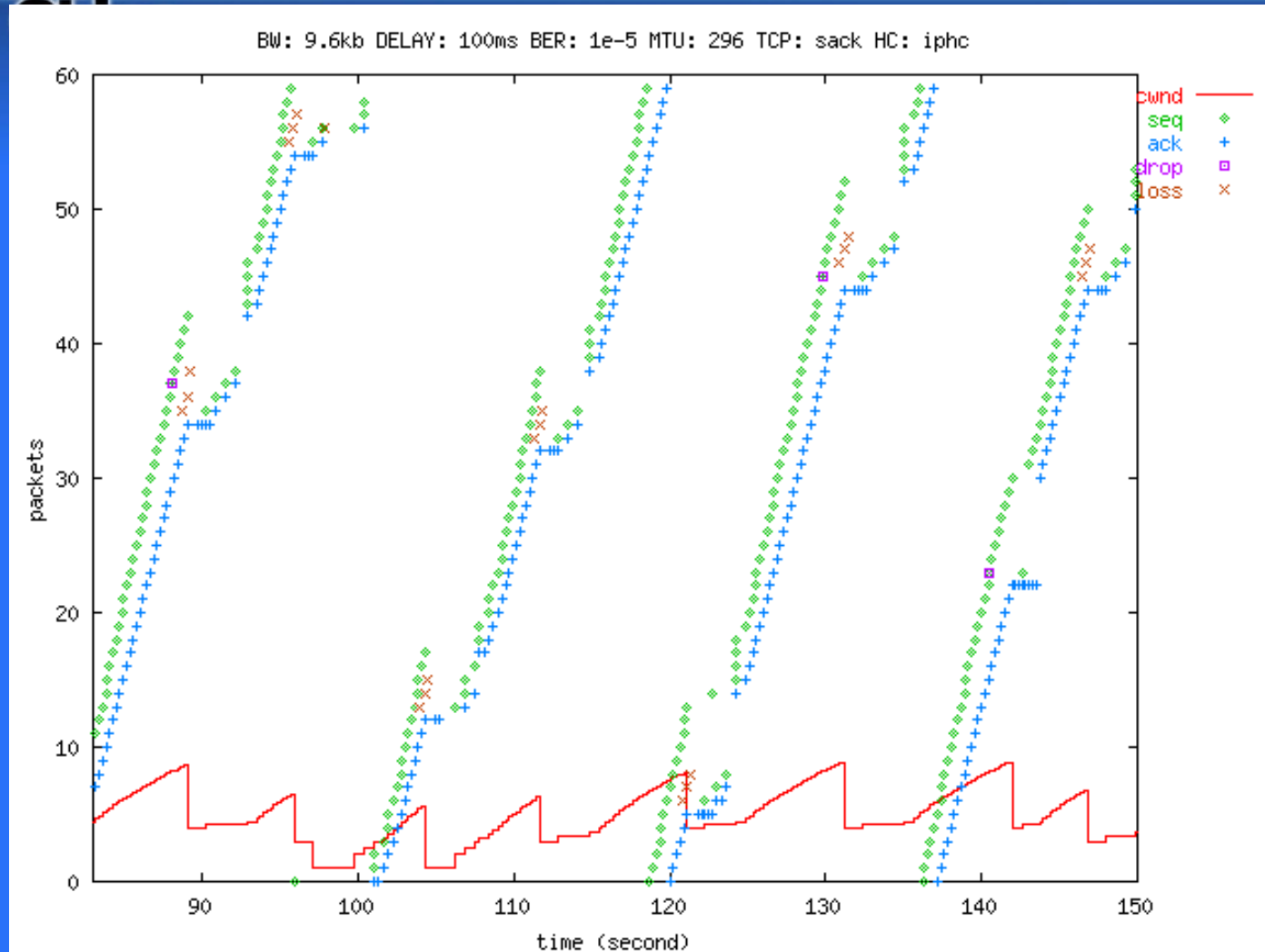
TCP behavior using RFC1144(VJHC)



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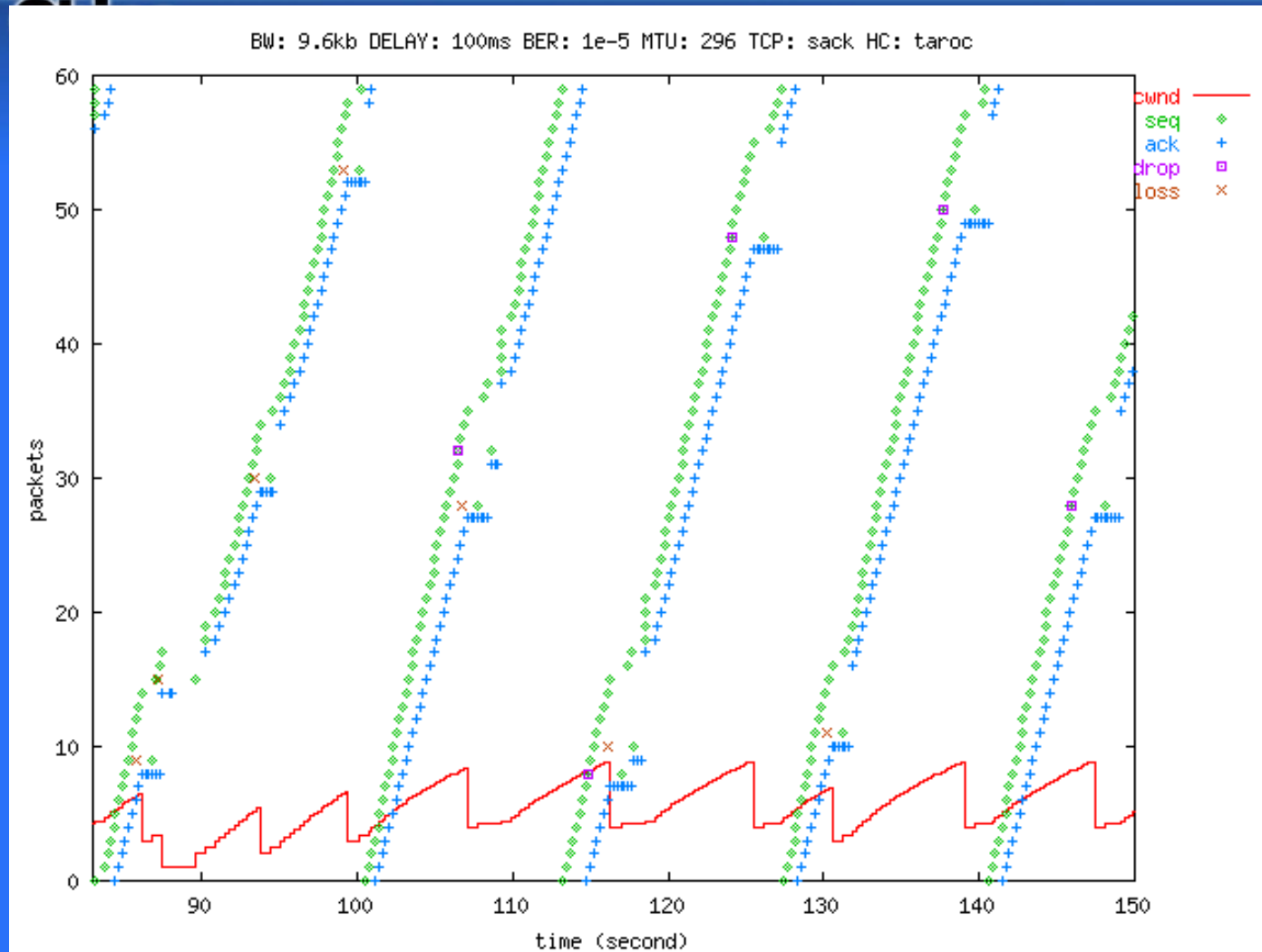
TCP behavior using RFC2507(IPHC)



TCP behavior using TAROC



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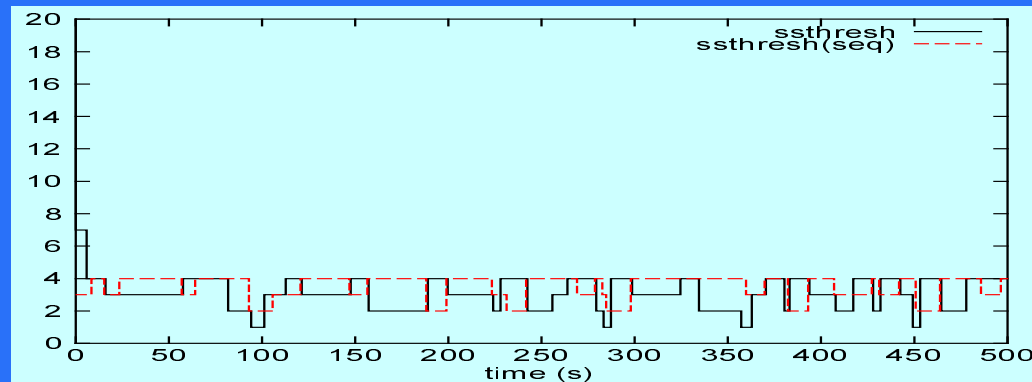
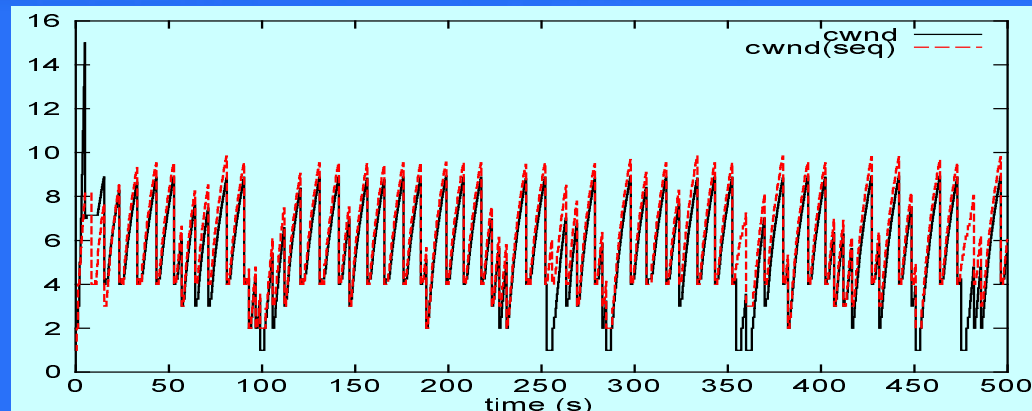


Accuracy of congestion window tracking



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- Tracking *cwnd* and *ssthresh* for seqno



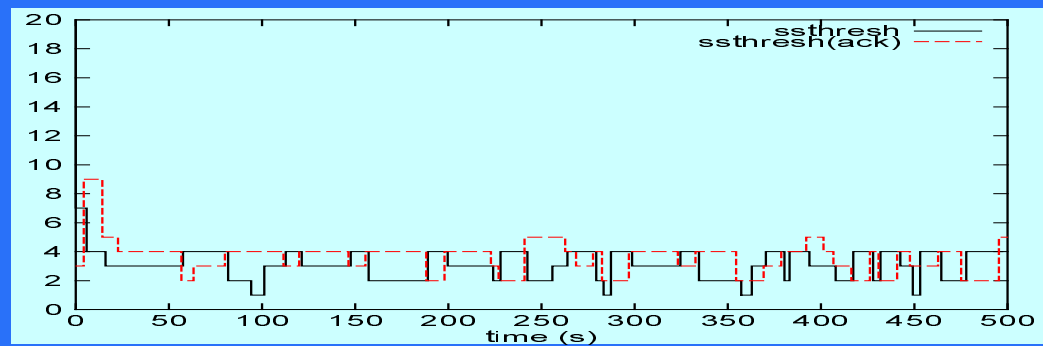
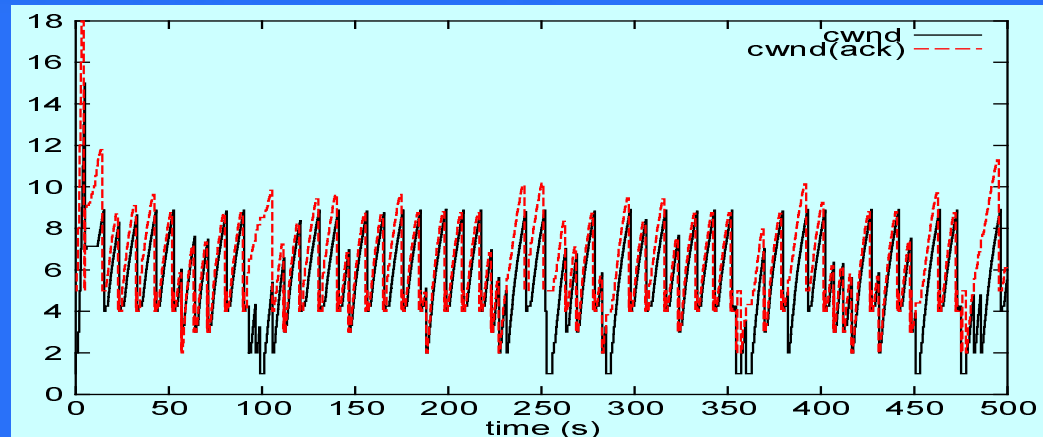
9.6kb-100ms-1e-5-296-sack

Accuracy of congestion window tracking



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- Tracking *cwnd* and *ssthresh* for ackno



9.6kb-100ms-1e-5-296-sack

Conclusions



- Robustness to
 - ✦ packet loss and misordering before the compressor
 - ✦ packet loss between compressor and decompressor
- Resilient to handover
 - ✦ seamlessly run across compressor/decompressor entity relocation
- Accurate TCP congestion window tracking
- High compression efficiency
 - ✦ average overhead at or slightly above 6 byte per packet, across a wide range of error rates
- Sufficient simulations demonstrate effective of TAROC



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Thanks for your attention!